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**FACULTY OF ENGINEERING AND TECHNOLOGY**

Distributed Systems (BCS-701)

LECTURE -16

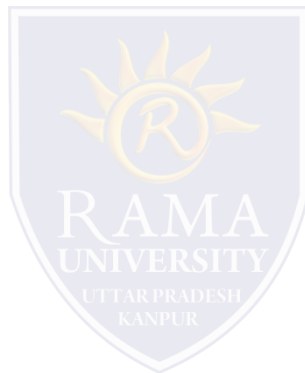
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# OUTLINE

- **Deadlock**
- **Fundamental Causes of Deadlock**
- **Deadlock Handling Strategies**
- 
- **MCQ**
- **Reference**



# DISTRIBUTED DEADLOCK DETECTION

## Deadlock

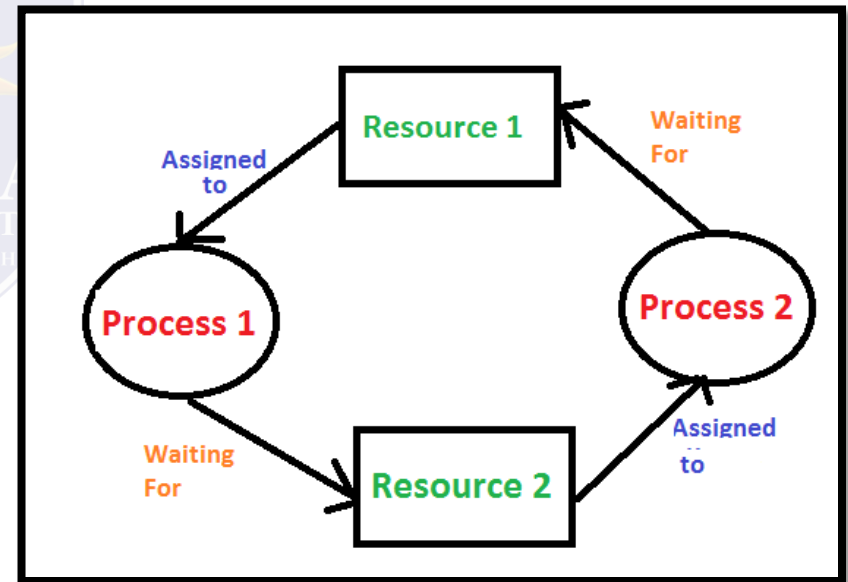
**Deadlock** is a situation where a set of processes are blocked because each process is holding a resource and waiting for another resource acquired by some other process.

### Resource deadlock:

each process requests resources held by another process in the set and it must receive all the requested resources before it can become unblocked.

### Communication deadlock:

each process is waiting for communication from another process, and will not communicate until it receives the communication for which it is waiting.



# DISTRIBUTED DEADLOCK DETECTION

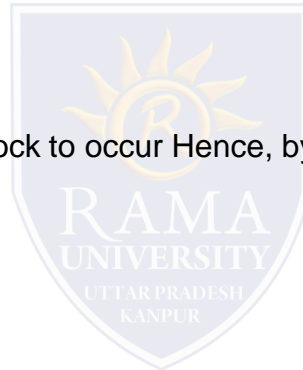
## Fundamental Causes of Deadlock

- exclusive access
- wait while holding
- no preemption
- circular wait

All these conditions are necessary for deadlock to occur Hence, by preventing any one of these we prevent deadlock.

## Deadlock Handling Strategies

- prevention
- avoidance
- detection



# DISTRIBUTED DEADLOCK DETECTION

## System Model

- The systems have only reusable resources.
- Processes are allowed only exclusive access to resources.
- There is only one copy of each resource.

## Distributed Deadlock Models

Based on WFG (not GRG)

- Nodes are processes
- Resources are located at a site, but may be held by processes at other sites
- Edge (P,Q) indicates P is blocked and waiting for Q to release some resource
- Deadlock exists if there is a directed cycle or knot.



# DISTRIBUTED DEADLOCK DETECTION

## Distributed Deadlock Handling Strategies

### Deadlock prevention

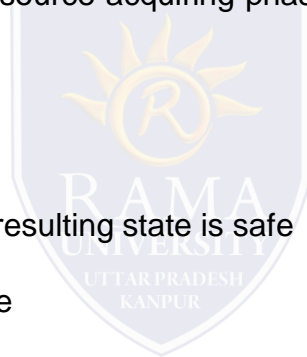
- All resource at once.
- Preventing a process from holding while waiting
- inefficient, can become deadlocked at resource acquiring phase, resource requirements are unpredictable -- not an efficient, universal solution.

### Deadlock avoidance

- A resource is granted to a process if the resulting state is safe
- Every site has to maintain the global state
- The checking for a safe state must be done with mutual exclusion
- The number of processes and resources in a distributed system is large
- Not a practical solution

### Deadlock detection

- Once deadlock, always deadlock -- detection won't be outdated
- deadlock detection can be preceed concurrently with normal activities
- This is the usual approach



# DISTRIBUTED DEADLOCK DETECTION

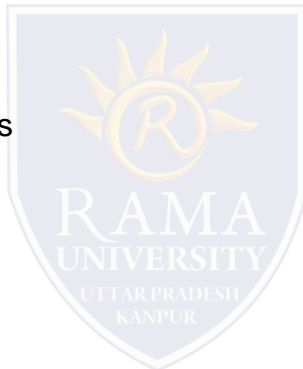
## Distributed Deadlock Detection Issues

### issues

- maintenance of WFG
- detection of cycles (or knots) in the WFG

### requirements

- progress = no undetected deadlocks
- safety = no false deadlocks



## Categorization of Methods

- centralized control
- distributed control
- hierarchical control

# DISTRIBUTED DEADLOCK DETECTION

## Classification of Distributed Detection Algorithms

### path-pushing

- ❑ path information transmitted, accumulated

### edge-chasing

- ❑ ``I'm waiting for you" probes are sent along edges
- ❑ single returned probe indicates a cycle

### diffusion

- ❑ ``Are you blocked?" probes are sent along all edges
- ❑ all queries returned indicates a cycle

### global state detection

- ❑ take and use snapshot of system state





# DISTRIBUTED DEADLOCK DETECTION

## Obermarck's Path-Pushing Algorithm

- ❑ designed for distributed database systems
- ❑ processes are called "transactions"  $T_1, T_2, \dots, T_n$
- ❑ there is special virtual node  $Ex$
- ❑ transactions are totally ordered

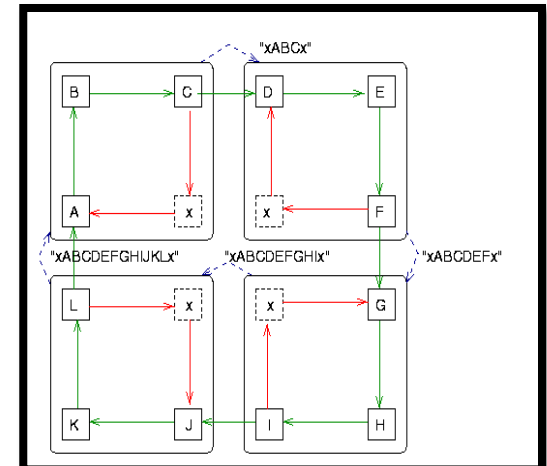
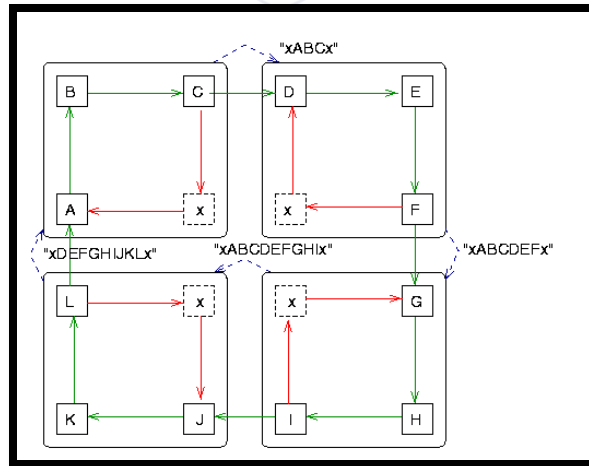
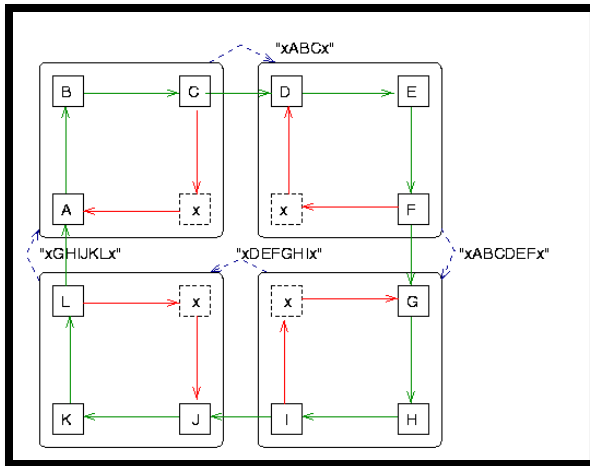
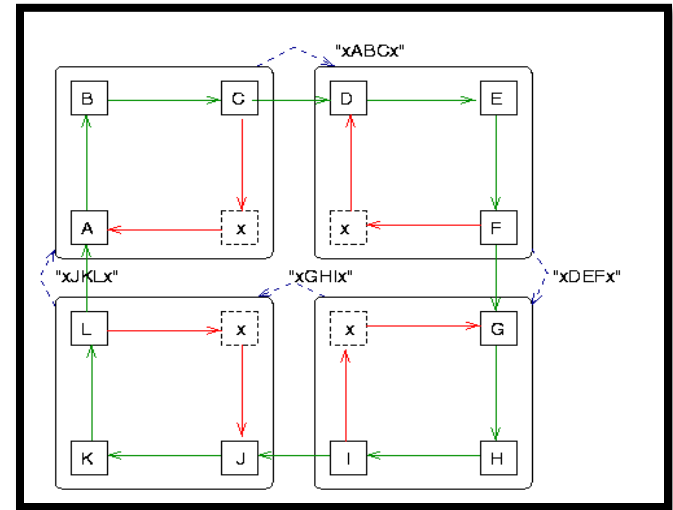
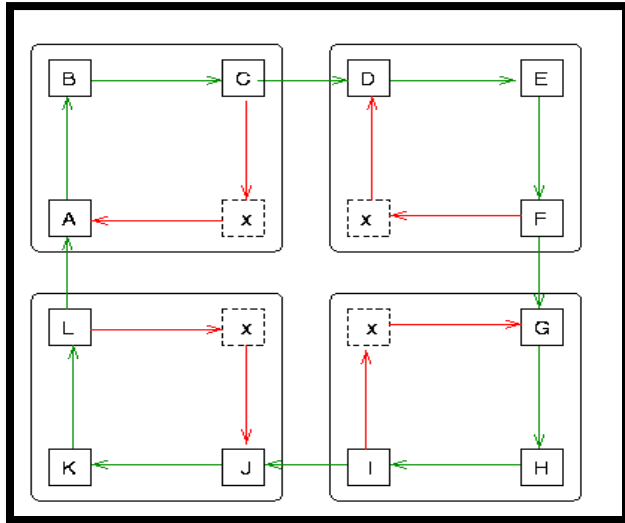
## Obermarck's Path-Pushing Algorithm

1. wait for info from previous iteration of Step 3
2. combine received info with local TWFG
  - ❑ detect all cycles
  - ❑ break local cycles
3. send each cycle including the node  $Ex$  to the external nodes it is waiting for
4. time-saver: only send path to other sites if last transaction is higher in lexical order than the first



# DISTRIBUTED DEADLOCK DETECTION

## Obermarck's Path-Pushing Algorithm



# DISTRIBUTED DEADLOCK DETECTION

## Problems with Obermarck's Path-Pushing Algorithm

- ❑ Detects false deadlocks, due to asynchronous snapshots at different sites.
- ❑ Message complexity? Message size? Detection delay?
- ❑ Exactly how are paths combined and checked?

## Obermarck's Path-Pushing Algorithm: Performance

- ❑  $O(n(n-1)/2)$  messages
- ❑  $O(n)$  message size
- ❑  $O(n)$  delay to detect deadlock



# DISTRIBUTED DEADLOCK DETECTION

## Chandy-Misra-Haas Edge-Chasing Algorithm

- ❑ for AND request model
- ❑ probe= (i,j,k) is sent for detection initiated by  $P_i$ , by site of  $P_j$  to site of  $P_k$
- ❑ deadlock is detected when a probe returns to its initiator

## Terminology

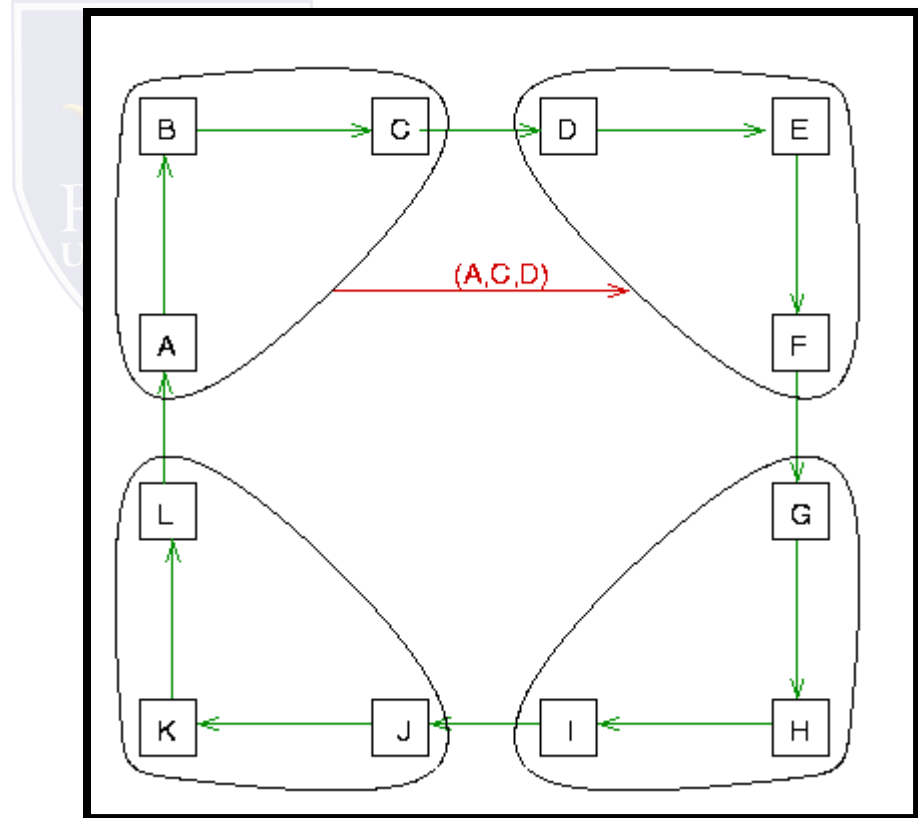
- ❑  $P_j$  is dependent on  $P_k$  if there is a sequence  $P_j, P_{i1}, P_{i2}, \dots, P_{im}, P_k$  such that each process except  $P_k$  is blocked and each process except the first holds a resource for which the previous process is waiting
- ❑  $P_j$  is locally dependent on  $P_k$  if it is dependent and both processes are at the same site
- ❑ array dependent  $i(j) = \text{true} \hat{=} P_i$  knows that  $P_j$  is dependent on it

# DISTRIBUTED DEADLOCK DETECTION

## Algorithm Initiation by $P_i$

if  $P_i$  is locally dependent on itself then declare a deadlock else send probe  $(i, j, k)$  to home site of  $P_k$  for each  $j, k$  such that all of the following hold

- ❑  $P_i$  is locally dependent on  $P_j$
- ❑  $P_j$  is waiting on  $P_k$
- ❑  $P_j$  and  $P_k$  are on different sites



# DISTRIBUTED DEADLOCK DETECTION

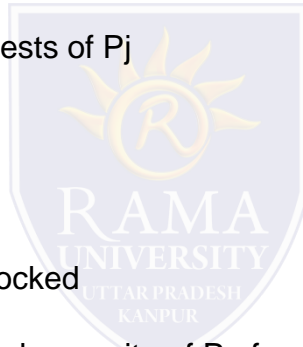
## Algorithm on receipt of probe (i,j,k)

check the following conditions

- $P_k$  is blocked
- Dependent  $k(i) = \text{false}$
- $P_k$  has not replied to all requests of  $P_j$

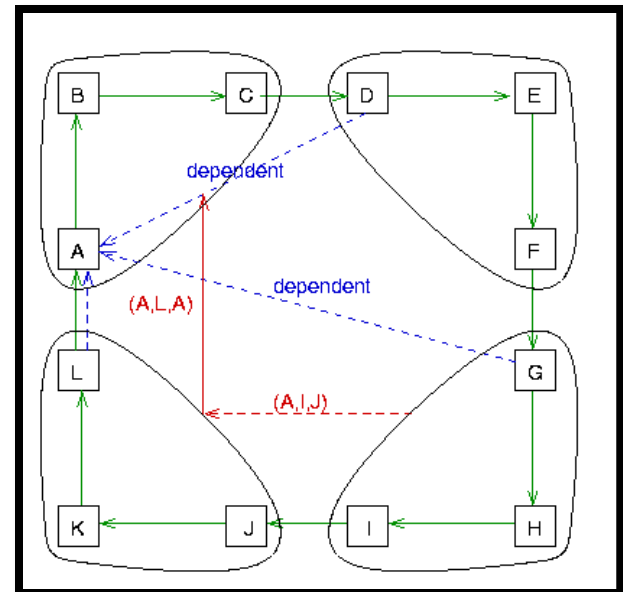
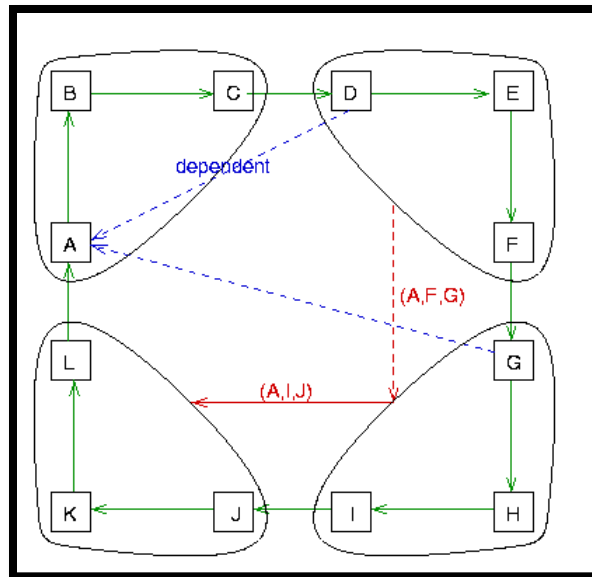
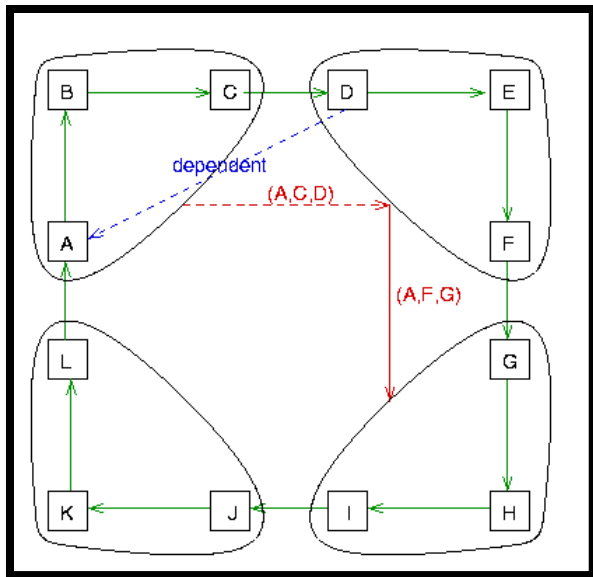
if these are all true, do the following

- set  $\text{dependent}_k(i) = \text{true}$
- if  $k=i$  declare that  $P_i$  is deadlocked
- else send probe (i,m,n) to the home site of  $P_n$  for every m and n such that the following all hold
  - $P_k$  is locally dependent on  $P_m$
  - $P_m$  is waiting on  $P_n$
  - $P_m$  and  $P_n$  are on different sites



# DISTRIBUTED DEADLOCK DETECTION

Algorithm on receipt of probe (i,j,k)

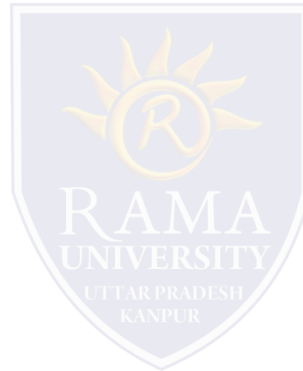


# DISTRIBUTED DEADLOCK DETECTION

## Algorithm on receipt of probe (i,j,k)

### Chandy Misra Haas Complexity

- ❑  $m(n-1)/2$  messages for m processes at n sites in the book, is this right?
- ❑ 3-word message length
- ❑  $O(n)$  delay to detect deadlock





# MCQ

1. What is the access point (AP) in a wireless LAN?

- a) device that allows wireless devices to connect to a wired network
- b) wireless devices itself
- c) both device that allows wireless devices to connect to a wired network and wireless devices itself
- d) all the nodes in the network

2. In wireless ad-hoc network \_\_\_\_\_

- a) access point is not required
- b) access point is must
- c) nodes are not required
- d) all nodes are access points

3. Which multiple access technique is used by IEEE 802.11 standard for wireless LAN?

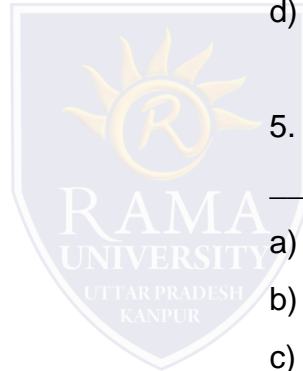
- a) CDMA
- b) CSMA/CA
- c) ALOHA
- d) CSMA/CD

4. In wireless distribution system \_\_\_\_\_

- a) multiple access point are inter-connected with each other
- b) there is no access point
- c) only one access point exists
- d) access points are not required

5. A wireless network interface controller can work in \_\_\_\_\_

- a) infrastructure mode
- b) ad-hoc mode
- c) both infrastructure mode and ad-hoc mode
- d) WDS mode



# REFERENCES

- ❑ <http://cs-www.cs.yale.edu/homes/aspnes/classes/465/notes.pdf>
- ❑ <https://www.geeksforgeeks.org/mutual-exclusion-in-distributed-system/>
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- ❑ <http://www.cs.fsu.edu/~xyuan/cop5611/lecture9.html>
- ❑ <http://www.cs.fsu.edu/~xyuan/cop5611/lecture10.html>

