



FACULTY OF ENGINEERING & TECHNOLOGY

Lecture : 01

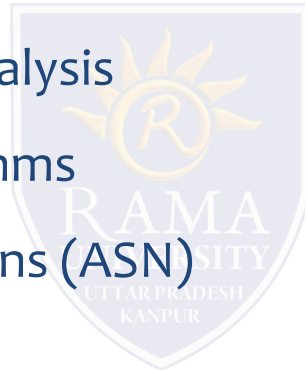
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□ Outline

- ❖ Introduction of Design and Analysis
- ❖ Analysis of Algorithms
- ❖ Need of Analysis
 - I. Methodology of Analysis
 - II. Behavior of Algorithms
 - III. Asymptotic Notations (ASN)



□ Introduction of Design and Analysis

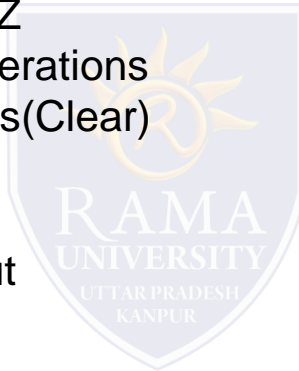
➤ ALGORITHMS

- ❖ Consists of Finite set of Steps to solve a problem

Ex: 1. $X = Y + Z$
 2. Read(a)
 3. for i 1 to n
 $X = Y + Z$

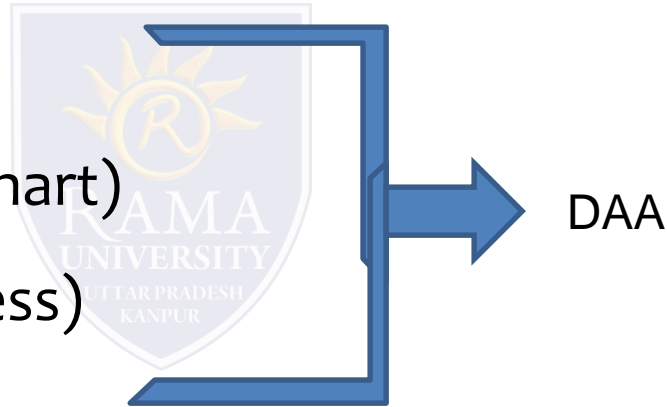
One or more operations

- ❖ I. definiteness(Clear)
- ❖ II. Effective
- ❖ May accept one or More Inputs
- ❖ Must produce at least on Output



□ Life Cycle Steps

1. Problem Define
2. Requirements Specification
3. Design(Logic)
4. Express(Algo/Flowchart)
5. Validation(correctness)
6. Analysis
7. Implementation
8. Testing & Debugging



□ Need of Algorithms

❖ Resource Comparison

1. Time
2. Space
3. Energy
4. Band width
5. Register



❖ Performance Comparison : Which also is best among all.

□ Methodology of Analysis

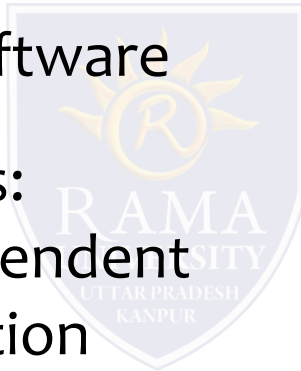
How much time taken by it

$$X=Y+Z$$

It depends on- Platform

- Hardware
- Software

1. Posterior Analysis:
 - I. Platform Dependent
 - II. Experimentation
2. Priori Analysis:
 - I. Platform Independent
 - II. Order of Magnitude
 - III. Operations: +,-,*



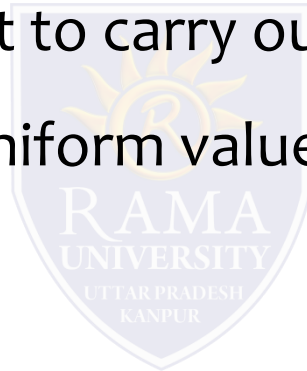
□ Need of Algorithms

❖ Posteriori Analysis

Advantages: Exact/ Real values of time & Space in units

Disadvantage: Difficult to carry out

Non- Uniform values gives.



❖ Oder of Magnitude: refers of frequency (No. of times) of the function/operations Involved in the steps/ statements.

□ Behavior of Algorithms

❖ Types of Analysis:

1. Worst case
2. Best case
3. Average Case

❖ Consider the Example

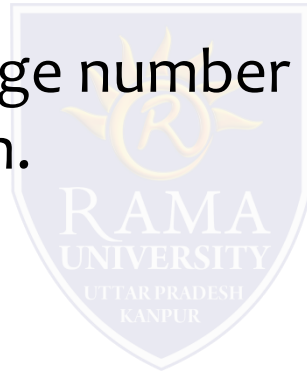
Given $\rightarrow (a_1, a_2, a_3, a_4 \dots a_n)$

I/P \rightarrow Inc. Order, I₁
Dec. Order, I₂
Random, I₃



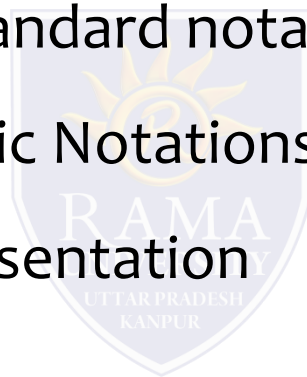
□ Behavior of Algorithms

- Worst-case: maximum number of steps taken on any instance of size n .
- Best-case: minimum number of steps taken on any instance of size n .
- Average case: average number of steps taken on any instance of size n .



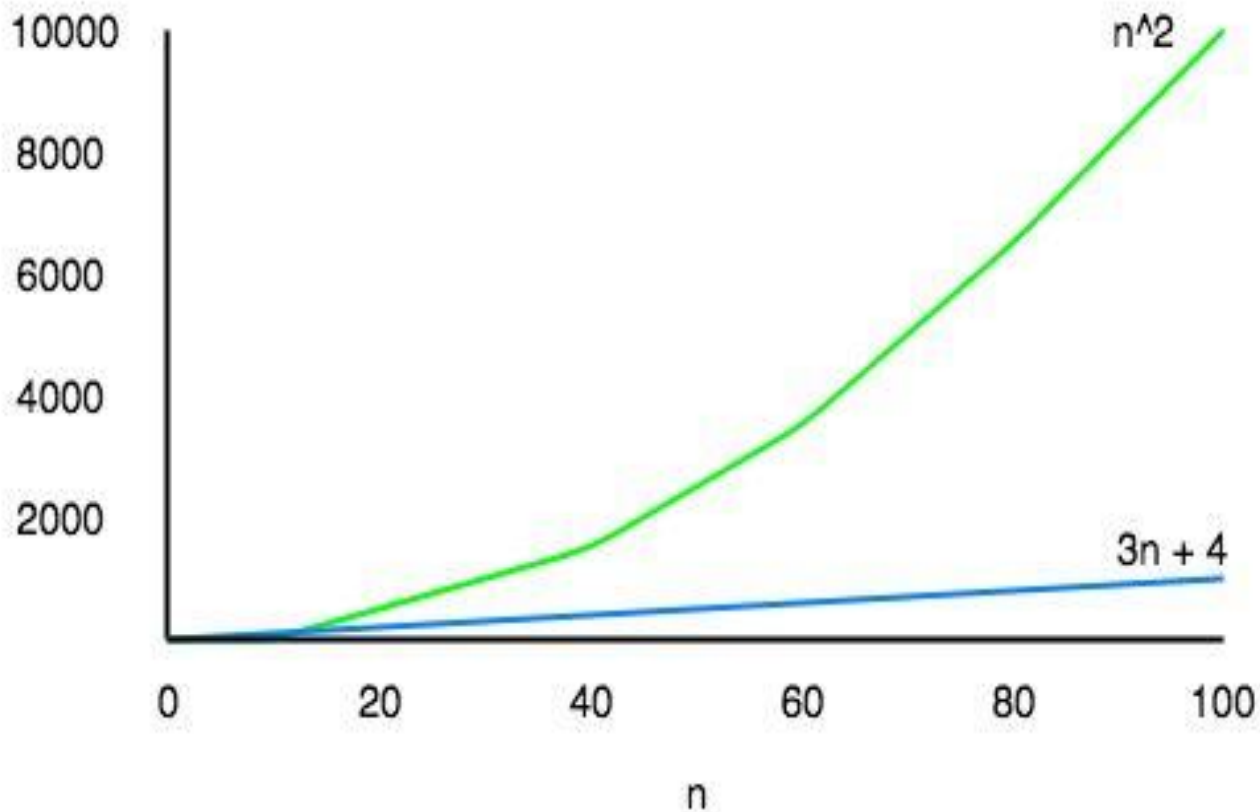
□ Asymptotic Notations

- We can never provide an exact number to define the time required and the space required by the algorithm.
- It is an express or standard notations of time and space, known as Asymptotic Notations.
- Mathematical Representation
- Bounds of Function
 1. Upper
 2. Lower
 3. Tight



□ Asymptotic Notations

- Time complexity of $T(n) = (n^2 + 3n + 4) \rightarrow$ quadratic equation
- For large values of n , the $3n + 4$ part will become insignificant compared to the n^2 part.



□ Asymptotic Notations

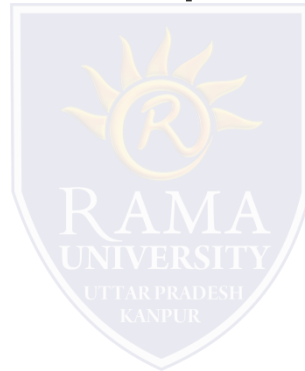
❖ Types of Asymptotic Notations

We use three types of asymptotic notations to represent the growth of any algorithm, as input increases:

1. Big Theta (Θ)

2. Big Oh (O)

3. Big Omega (Ω)



□ Big Oh(O)

- Given functions $f(n)$ and $g(n)$, we say that $f(n)$ is $O(g(n))$ if there are positive constants c and n_0 such that

$$f(n) \leq cg(n) \text{ for } n \geq n_0$$

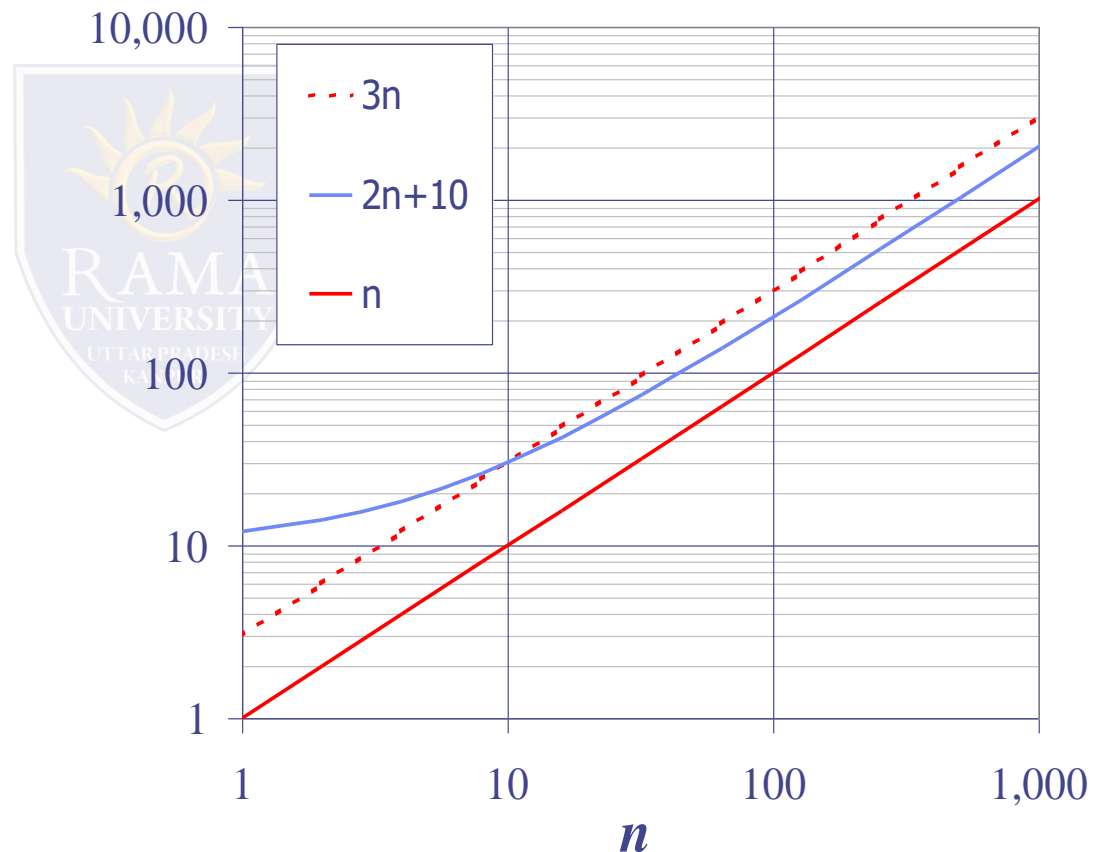
- Example: $2n + 10$ is $O(n)$

- $2n + 10 \leq cn$

- $(c - 2)n \geq 10$

- $n \geq 10/(c - 2)$

- Pick $c = 3$ and $n_0 = 10$



□ Big Omega (Ω) & Big-Theta

→big-Omega

- ❖ $f(n)$ is $\Omega(g(n))$ if there is a constant $c > 0$
- ❖ and an integer constant $n_0 \geq 1$ such that
- ❖ $f(n) \geq c \cdot g(n)$ for $n \geq n_0$

→big-Theta

- $f(n)$ is $\Theta(g(n))$ if there are constants $c' > 0$ and $c'' > 0$ and an integer constant $n_0 \geq 1$ such that $c' \cdot g(n) \leq f(n) \leq c'' \cdot g(n)$ for $n \geq n_0$



→ Big-Oh

- $f(n)$ is $O(g(n))$ if $f(n)$ is asymptotically less than or equal to $g(n)$

→ big-Omega

- $f(n)$ is $\Omega(g(n))$ if $f(n)$ is asymptotically greater than or equal to $g(n)$

→ big-Theta

- $f(n)$ is $\Theta(g(n))$ if $f(n)$ is asymptotically equal to $g(n)$



Small/Little Notation:

- Small- Oh: Proper Upper Bond
- Small-Omega: Proper Lower Bond

□ Some common asymptotic notations –

constant	–	$O(1)$
logarithmic	–	$O(\log n)$
linear	–	$O(n)$
$n \log n$	–	$O(n \log n)$
quadratic	–	$O(n^2)$
cubic	–	$O(n^3)$
polynomial	–	$n^{O(1)}$
exponential	–	$2^{O(n)}$

□ Practice

Q: What is time complexity of fun()?

```
int fun(int n)
{
    int count = 0;
    for (int i = n; i > 0; i /= 2)
        for (int j = 0; j < i; j++)
            count += 1;
    return count;
}
```



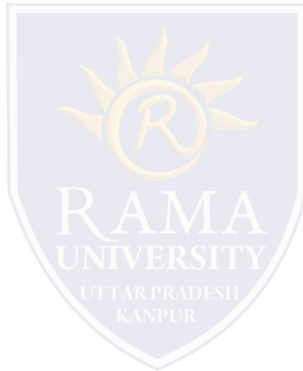
Options:

- A. $O(n^2)$
- B. $O(n \log n)$
- C. $O(n)$
- D. $O(n \log n \log n)$

□ Practice

Q: What is the time complexity of fun()?

```
int fun(int n)
{
    int count = 0;
    for (int i = 0; i < n; i++)
        for (int j = i; j > 0; j--)
            count = count + 1;
    return count;
}
```



Options:

- A. Theta (n)
- B. Theta (n²)
- C. Theta (n*Logn)
- D. Theta (nLognLogn)